

# Subgraph Deletion of 4-Regular Graphs and Their Genus Ranges

Hayden Hunter

University of South Florida

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# Definition: Double Occurrence Word

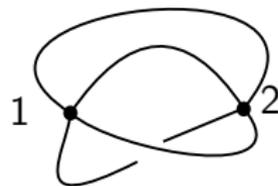
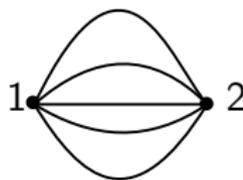
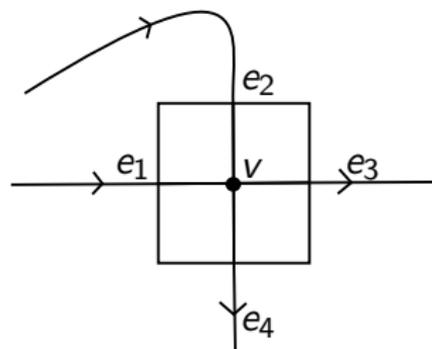
## Definition

Let  $A$  be an alphabet. A *double occurrence word* over  $A$  is a word which contains each symbol of  $A$  exactly 0 or 2 times. We denote the set of double occurrence words as  $A_{DOW}$ .

Example: Let  $A = \mathbb{N}$ . Then  $121323 \in A_{DOW}$

# Definition: Rigid Vertex

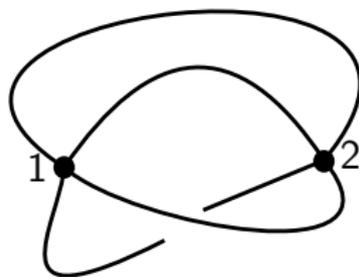
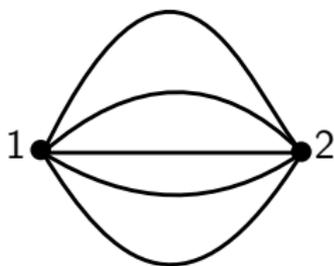
- 1 Rotation System
- 2 Sharp Corners are not permitted



# Definition: Assembly Graph

## Definition: Assembly Graph

Let  $\Gamma$  be a graph where each vertex is a rigid vertex of degree 4 or 1. Then we call  $\Gamma$  an *assembly graph*.



# Definition: Rigid Vertex and Assembly Graph

The graph below is the assembly graph of 121323. For  $\Gamma = (V(\Gamma), E(\Gamma))$  we have that

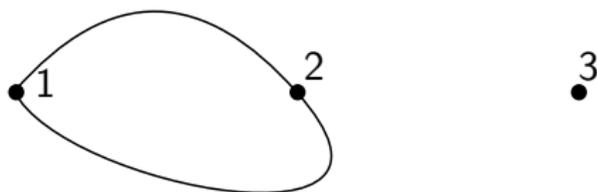
$$1 \rightarrow 2$$



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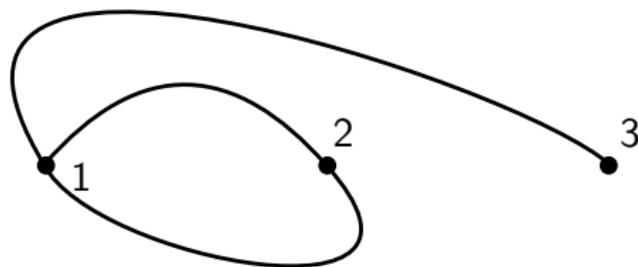
$$1 \rightarrow 2 \rightarrow 1$$



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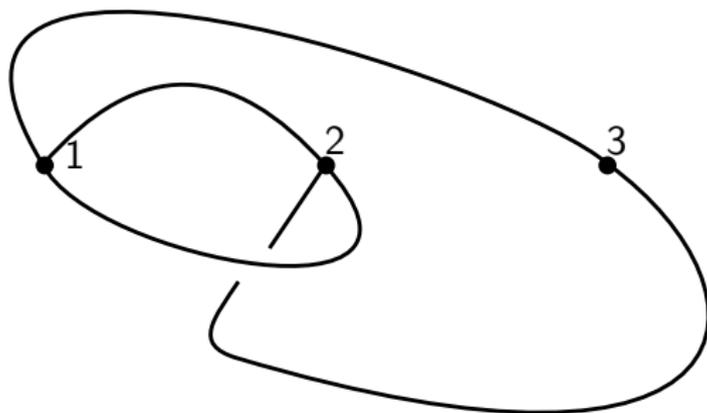
$$1 \rightarrow 2 \rightarrow 1 \rightarrow 3$$



# Definition: Rigid Vertex and Assembly Graph

The graph below is the assembly graph of 121323. For  $\Gamma = (V(\Gamma), E(\Gamma))$  we have that

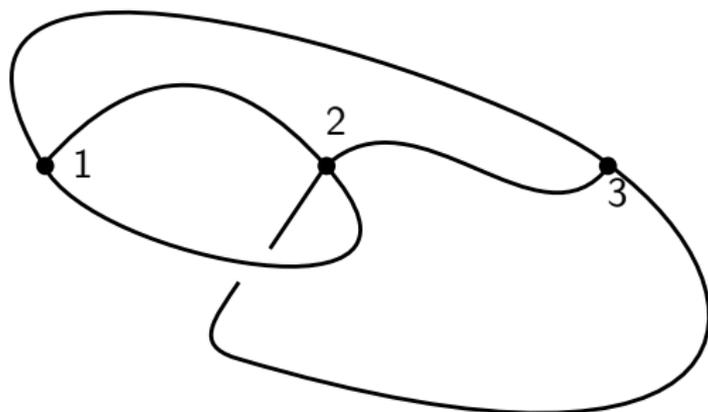
$$1 \rightarrow 2 \rightarrow 1 \rightarrow 3 \rightarrow 2$$



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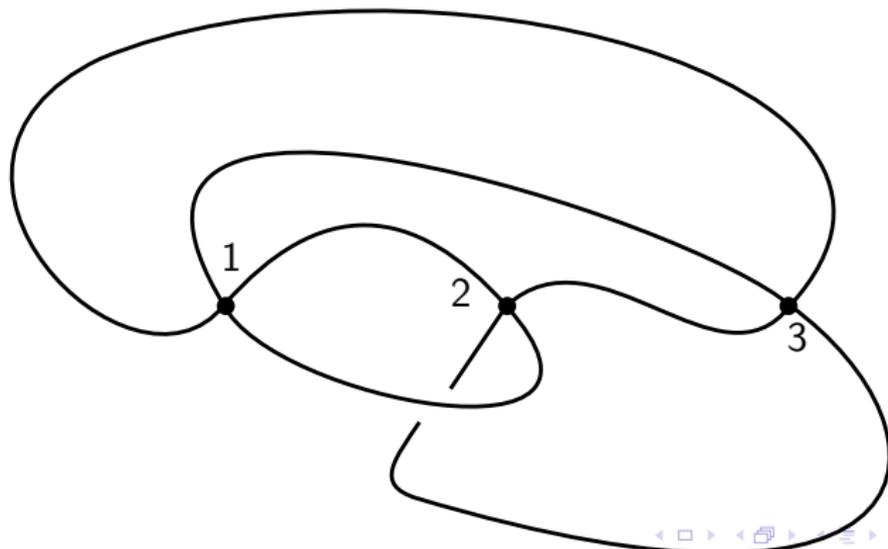
$$1 \rightarrow 2 \rightarrow 1 \rightarrow 3 \rightarrow 2 \rightarrow 3$$



# Definition: Rigid Vertex and Assembly Graph

The graph below is the assembly graph of 121323. For  $\Gamma = (V(\Gamma), E(\Gamma))$  we have that

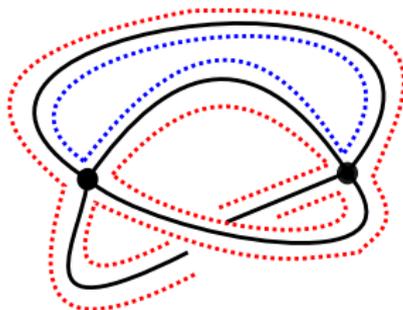
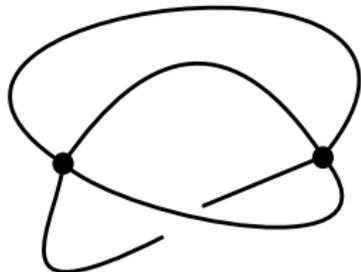
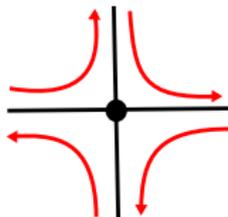
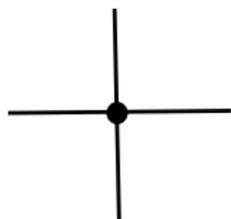
$$1 \rightarrow 2 \rightarrow 1 \rightarrow 3 \rightarrow 2 \rightarrow 3 \rightarrow 1$$



# Graphs and Their Genus

- Ribbon graphs and boundary components
- The Euler Characteristic of these ribbon graphs
- The change of a boundary connection at a vertex

# Ribbon Graphs



# Euler Characteristic of Ribbon Graphs

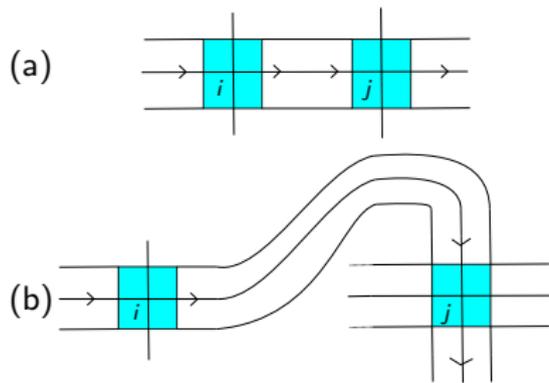
## Definition

The integer  $g$  that represents the number of handles a topological space has is the *genus* of that topological space

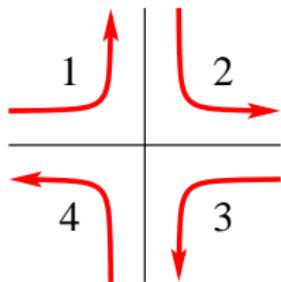
- The genus of an assembly graph  $\Gamma$  is the minimum number of handles of the topological space  $\Gamma$  can be embedded into.
- Letting  $b(\Gamma)$  be the number of boundary components, and considering that each assembly graph with  $n$  vertices has  $2n$  edges,

$$\chi = n - 2n + b(\Gamma) \implies g(\Gamma) = \frac{1}{2}(2 + n - b(\Gamma))$$

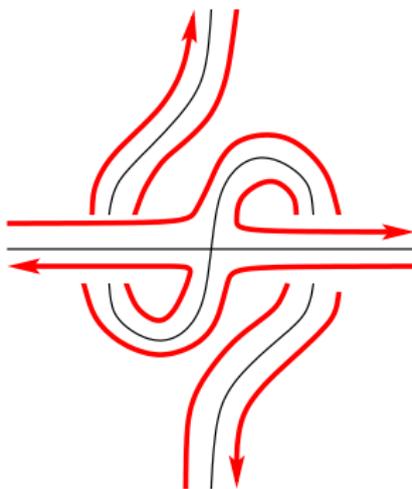
For each vertex there are two possible ways to construct our ribbon graph at the second occurrence.



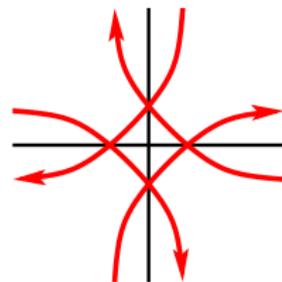
Thus we have that for an assembly graph  $\Gamma$  with  $n$  vertices, that there exists  $2^n$  boundary connections and thus  $2^n$  different graphical embeddings.



(A)



(B)



(C)

## Definition: Genus Range

Letting  $a$  be the minimum genus of  $\Gamma$  and  $b$  being the maximum genus of  $\Gamma$ , we say that the *genus range* of  $\Gamma$  is  $[a, b]$ .

- The minimum genus will be the minimum number of handles of the topological space that  $\Gamma$  is being embedded onto.
- The maximum genus will be the maximum number of handles of the topological space that  $\Gamma$  is being embedded onto.

# Repeat and Return Insertions

Let  $u = u_1 u_2 \dots u_n$ . Then we say that the reverse of  $u$  denoted as  $u^R = u_n u_{n-1} \dots u_2 u_1$ .

## Definition: Repeat and Return Insertions

Let  $w = xyz \in A$ . Let  $u$  be a single occurrence word where  $w \cap u = \{\varepsilon\}$  where  $|\varepsilon| = 0$ . Then  $\mathcal{T}(u)$  acts on  $w$  so that

$$w \star \mathcal{T}(u) = xuyu'z$$

where

$$u' = \begin{cases} u & \mathcal{T} = \rho \\ u^R & \mathcal{T} = \tau \end{cases}$$

We call  $\rho$  the repeat insertion and  $\tau$  the return insertion.

## Example: Repeat and Return Insertion

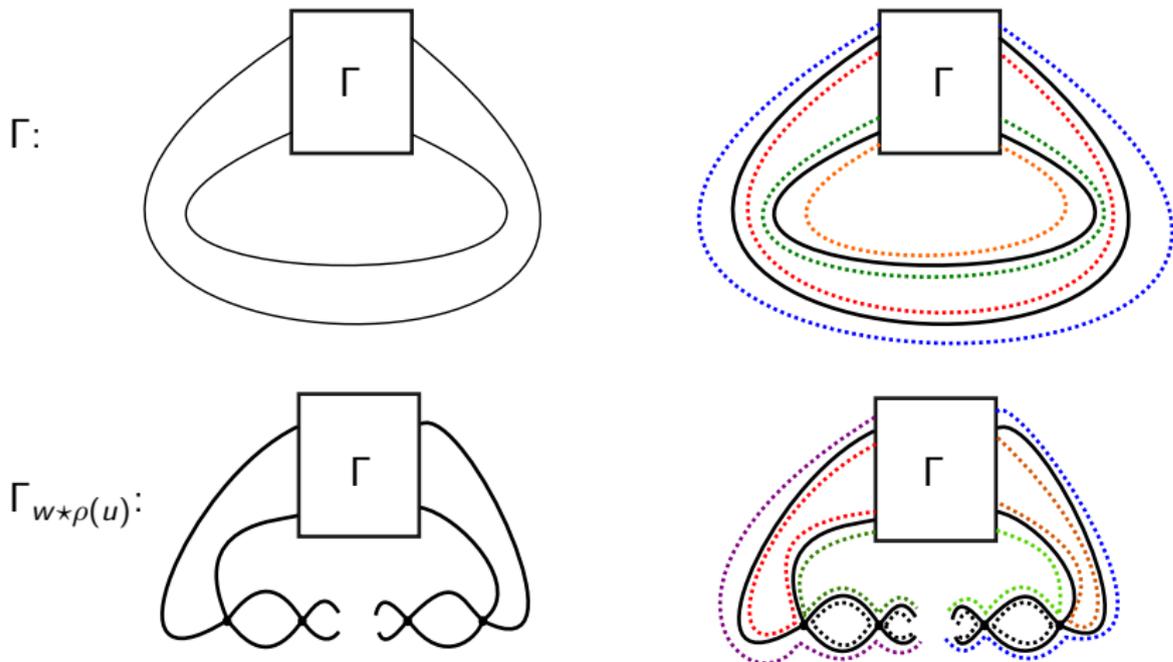
Let  $w = 121323$  and let  $u = 45$ . Then we have that

$$w \star \rho(u) = 1452134523$$

$$w \star \tau(u) = 1452135423$$

# Graphical Representation of a Repeat Insertion

Let  $\Gamma_w$  be the assembly graph for some double occurrence word  $w = xyz \in A_{DOW}$ . Let  $u$  be a single occurrence word of length  $m$ . Then we can represent  $\Gamma_w$  and  $\Gamma_{w \star \rho(u)}$  as



# External Connection Graphs

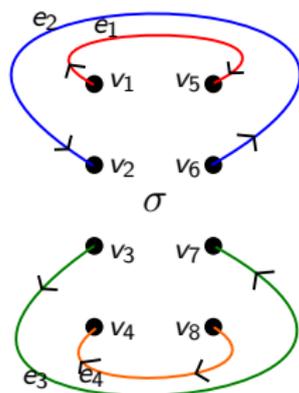
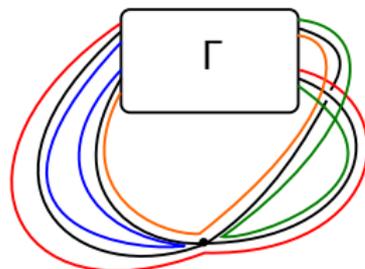
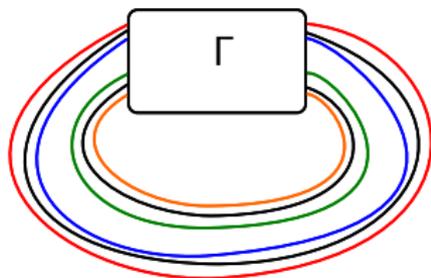


Fig 3:  $G$

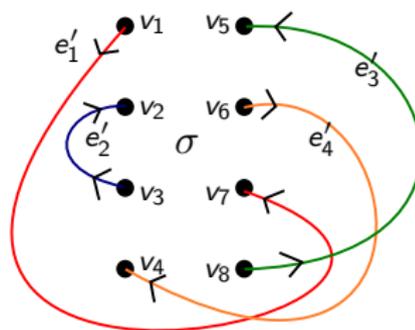


Fig 4:  $G'$

## Definition: External Connection Graph

**Definition 1.** Let  $w \in A_{DOW}$  for some alphabet  $A$ . Let  $G = (V, E)$  be a graph where  $V(G) = \{v_1, \dots, v_8\}$ ,  $\sigma' : \{v_1, v_3, v_6, v_8\} \rightarrow \{v_2, v_4, v_5, v_7\}$  and  $\sigma : \{v_2, v_4, v_5, v_7\} \rightarrow \{v_1, v_3, v_6, v_8\}$  are bijective functions, and  $E(G) = E_\sigma \dot{\cup} E_{\text{ext}}$  where

$$E_{\text{ext}} = \{(v_i, \sigma'(v_i)) \mid i = 1, 3, 6, 8\}$$

$$E_\sigma = \{(v_i, \sigma(v_i)) \mid i = 2, 4, 5, 7\}$$

We call  $G$  the **exterior connection graph** of  $\Gamma_{w \star \rho(u, i, j)}$ .

# Theorem: Single Insertion of $w$

## Theorem 1

**Theorem 1.** *Let  $A = \mathbb{N} \setminus \{1\}$  be an alphabet and  $w \in A_{DOW}$ . Let  $\Gamma = \Gamma_w$  be the assembly graph for  $w$  whose genus range is  $[a, b]$ . Then the assembly graph  $\Gamma' = \Gamma_{w \star \rho(1)}$  for the double occurrence word  $w \star \rho(1)$  has a genus range of  $[a + \epsilon, b + \epsilon']$  where  $\epsilon, \epsilon' \in \{-1, 0, 1, 2\}$*

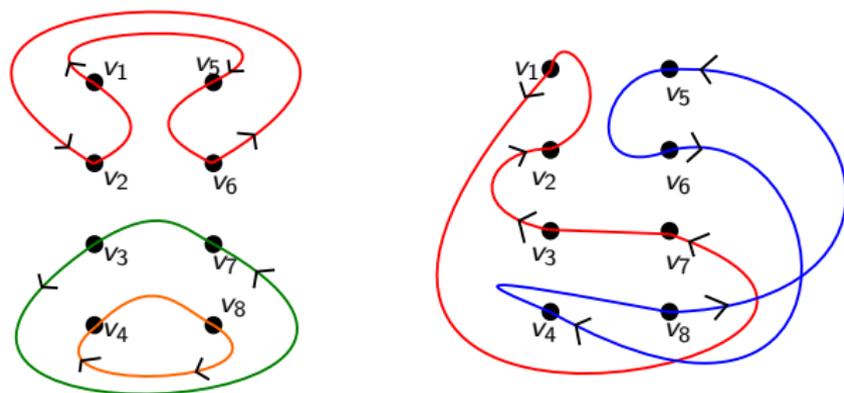
$M_G$	$e_1$	$e_2$	$e_3$	$e_4$	$e_5$	$e_6$	$e_7$	$e_8$
$v_1$	+1	0	0	0	-1	0	0	0
$v_2$	0	-1	0	0	$\zeta_{25}$	$\zeta_{26}$	$\zeta_{27}$	$\zeta_{28}$
$v_3$	0	0	+1	0	0	-1	0	0
$v_4$	0	0	0	-1	$\zeta_{45}$	$\zeta_{46}$	$\zeta_{47}$	$\zeta_{48}$
$v_5$	-1	0	0	0	$\zeta_{55}$	$\zeta_{56}$	$\zeta_{57}$	$\zeta_{58}$
$v_6$	0	+1	0	0	0	0	-1	0
$v_7$	0	0	-1	0	$\zeta_{75}$	$\zeta_{76}$	$\zeta_{77}$	$\zeta_{78}$
$v_8$	0	0	0	+1	0	0	0	-1

↓

$M'_G$	$e_1$	$e_2$	$e_3$	$e_4$	$e_5$	$e_6$	$e_7$	$e_8$
$v_1$	+1	0	0	0	-1	0	0	0
$v_2$	0	-1	0	0	$\zeta_{25}$	$\zeta_{26}$	$\zeta_{27}$	$\zeta_{28}$
$v_3$	0	+1	0	0	0	-1	0	0
$v_4$	0	0	-1	0	$\zeta_{45}$	$\zeta_{46}$	$\zeta_{47}$	$\zeta_{48}$
$v_5$	0	0	0	-1	$\zeta_{55}$	$\zeta_{56}$	$\zeta_{57}$	$\zeta_{58}$
$v_6$	0	0	+1	0	0	0	-1	0
$v_7$	-1	0	0	0	$\zeta_{75}$	$\zeta_{76}$	$\zeta_{77}$	$\zeta_{78}$
$v_8$	0	0	0	+1	0	0	0	-1

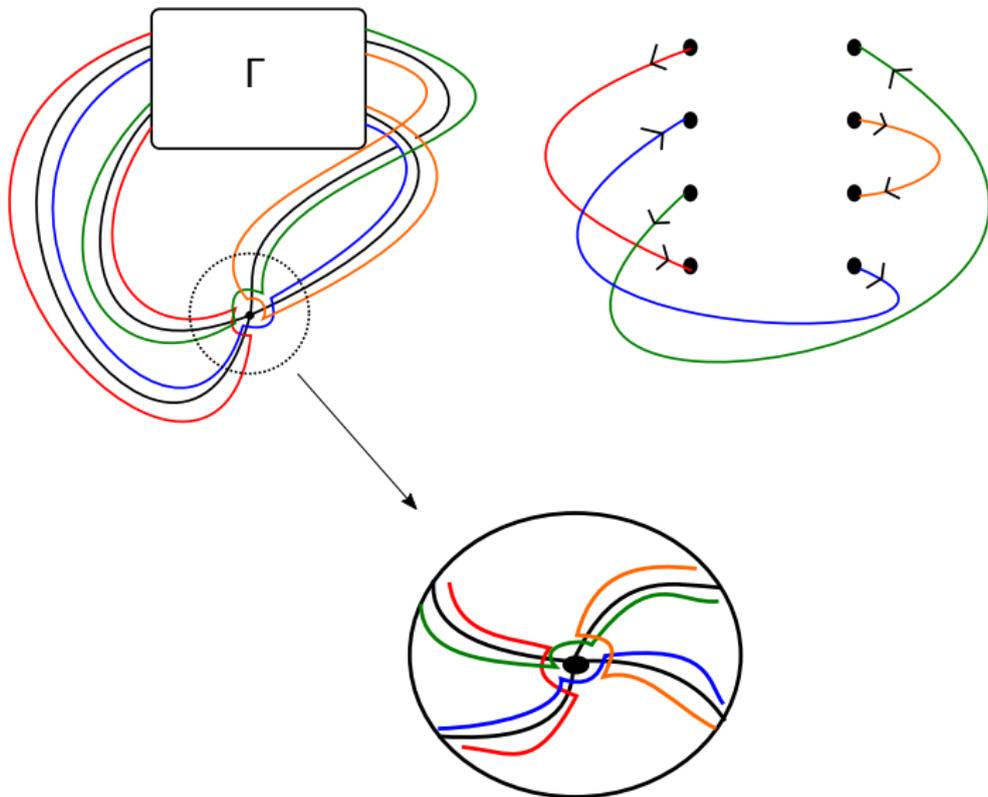
## Example

Let  $\sigma((v_1, v_3, v_6, v_8)) = (v_2, v_7, v_5, v_4)$ . Then the external connection graphs for  $\Gamma_w$  and  $\Gamma_{w \star xyz}$  are



Permutation Type	$ b(\Gamma) $	$ b(\Gamma') $	$ b(\Gamma')  -  b(\Gamma) $
$(v_2, v_4, v_5, v_7)$	2	1	-1
$(v_2, v_4, v_7, v_5)$	1	2	1
$(v_2, v_5, v_4, v_7)$	1	2	1
$(v_2, v_5, v_7, v_4)$	2	1	-1
$(v_2, v_7, v_4, v_5)$	2	3	1
$(v_2, v_7, v_5, v_4)$	3	2	-1
$(v_4, v_2, v_5, v_7)$	1	2	1
$(v_4, v_2, v_7, v_5)$	2	3	1
$(v_4, v_5, v_2, v_7)$	2	1	-1
$(v_4, v_5, v_7, v_2)$	1	2	1
$(v_4, v_7, v_2, v_5)$	3	2	-1
$(v_4, v_7, v_5, v_2)$	2	1	-1
$(v_5, v_2, v_4, v_7)$	2	3	1
$(v_5, v_2, v_7, v_4)$	3	2	-1
$(v_5, v_4, v_2, v_7)$	3	2	-1
$(v_5, v_4, v_7, v_2)$	2	1	-1
$(v_5, v_7, v_2, v_4)$	4	1	-3
$(v_5, v_7, v_4, v_2)$	3	2	-1
$(v_7, v_2, v_4, v_5)$	1	4	3
$(v_7, v_2, v_5, v_4)$	2	3	1
$(v_7, v_4, v_2, v_5)$	2	3	1
$(v_7, v_4, v_5, v_2)$	1	2	1
$(v_7, v_5, v_2, v_4)$	3	2	-1
$(v_7, v_5, v_4, v_2)$	2	3	1

We then consider what occurs when we change the boundary at the inserted vertex.  $G''$  in this context is the external connection graph for  $\Gamma$  with a boundary connection change at the inserted vertex.



Permutation Type	$ b(\Gamma) $	$ b'(\Gamma') $	$ b'(\Gamma')  -  b(\Gamma) $
$(v_2, v_4, v_5, v_7)$	2	1	-1
$(v_2, v_4, v_7, v_5)$	1	2	1
$(v_2, v_5, v_4, v_7)$	1	2	1
$(v_2, v_5, v_7, v_4)$	2	3	1
$(v_2, v_7, v_4, v_5)$	2	1	-1
$(v_2, v_7, v_5, v_4)$	3	2	-1
$(v_4, v_2, v_5, v_7)$	1	2	1
$(v_4, v_2, v_7, v_5)$	2	3	1
$(v_4, v_5, v_2, v_7)$	2	3	1
$(v_4, v_5, v_7, v_2)$	1	4	3
$(v_4, v_7, v_2, v_5)$	3	2	-1
$(v_4, v_7, v_5, v_2)$	2	3	1
$(v_5, v_2, v_4, v_7)$	2	1	-1
$(v_5, v_2, v_7, v_4)$	3	2	-1
$(v_5, v_4, v_2, v_7)$	3	2	-1
$(v_5, v_4, v_7, v_2)$	2	3	1
$(v_5, v_7, v_2, v_4)$	4	1	-3
$(v_5, v_7, v_4, v_2)$	3	2	-1
$(v_7, v_2, v_4, v_5)$	1	2	1
$(v_7, v_2, v_5, v_4)$	2	1	-1
$(v_7, v_4, v_2, v_5)$	2	1	-1
$(v_7, v_4, v_5, v_2)$	1	2	1
$(v_7, v_5, v_2, v_4)$	3	2	-1
$(v_7, v_5, v_4, v_2)$	2	3	1

$$g(\Gamma') = \frac{1}{2}(2 + (n + 1) - b(\Gamma')) = \frac{1}{2}(2 + n - b(\Gamma')) + \frac{1}{2}(1 - \epsilon)$$

where  $\epsilon = -3, -1, 1, 3$  so that

$$g(\Gamma') = g(\Gamma) + \delta$$

where  $\delta = -1, 0, 1, 2$ .

## Theorem 2

**Theorem 2.** *Let  $A$  be an alphabet and  $u, v, w \in A$  so that  $uvw \in A_{DOW}$ . Let  $\Gamma$  be the assembly graph for  $uvw$  whose genus range is  $[a, b]$ . Then the assembly graph  $\Gamma'$  for the double occurrence word  $u12v12w$  has a genus range of  $[a + \epsilon, b + \epsilon']$  where  $\epsilon, \epsilon' \in \{0, 1, 2\}$*

## Theorem 3

**Theorem 3** Let  $A$  be an alphabet  $w \in A_{DOW}$ . Let  $\Gamma = \Gamma_w$  be the assembly graph for  $w$  whose genus range is  $[a, b]$ . Let  $v_{odd}, v_{even}, u \in A_{SOW}$  where  $|v_{odd}| = 3, |v_{even}| = 2$ , and  $|u| = n$ .

$$g(\Gamma_{w \star \mathcal{T}(u, i, j)}) = \begin{cases} g(\Gamma_{w \star \mathcal{T}(v_{odd}, i, j)}) & \text{If } n \equiv \text{odd} \\ g(\Gamma_{w \star \mathcal{T}(v_{even}, i, j)}) & \text{If } n \equiv \text{even} \end{cases}$$

## Example

Consider the *DOW*

$$12145673234567 = 121323 \star \rho(4567)$$

## Example

Consider the *DOW*

$$12145673234567 = 121323 \star \rho(4567)$$

We have that  $g(121323 \star \rho(4567)) = g(121323 \star \rho(45)) = [2, 3]$

# Conclusion

- Construction of a lower bound
- Classes of assembly graphs after subgraph deletion

- D. Buck, E. Dolzhenko, N. Jonoska, M. Saito, K. Valencia. Genus Ranges of 4-Regular Rigid Vertex Graphs. *The Electronic Journal of Combinatorics*. pg (1-10,15)
- N. Jonoska, L. Nabergall, M. Saito. Patterns and Distances in Words Related to DNA Rearrangement. *Fundamenta Informaticae*. pg(1003 - 1007)
- D. Cruz, M. Ferrari, N. Jonoska, L. Nabergall, M. Saito. Transformations on Double Occurrence Words Motivated by DNA Rearrangement. arXiv:1811.11739 . pg(5,6)

Thank you!

